Describing Admissible Rules

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April 28^{th} 2014

Disjunction Property

$$\frac{A \vee B \text{ derivable}}{A \text{ derivable or } B \text{ derivable}}$$

$$\frac{\vdash A \vee B}{\vdash A \text{ or } \vdash B}$$

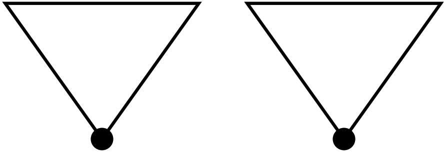
syntax

$$\frac{\vdash A \vee B}{\vdash A \text{ or } \vdash B}$$

semantics

syntax

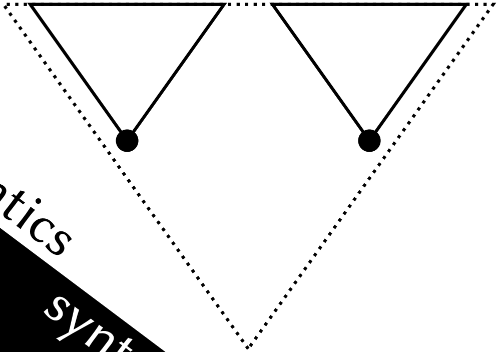
$$\frac{\vdash A \vee B}{\vdash A \text{ or } \vdash B}$$



semantics

syntax

$$\frac{\vdash A \vee B}{\vdash A \text{ or } \vdash B}$$



semantics

syntax

$$\frac{\vdash A \vee B}{\vdash A \text{ or } \vdash B}$$

Overview

Overview

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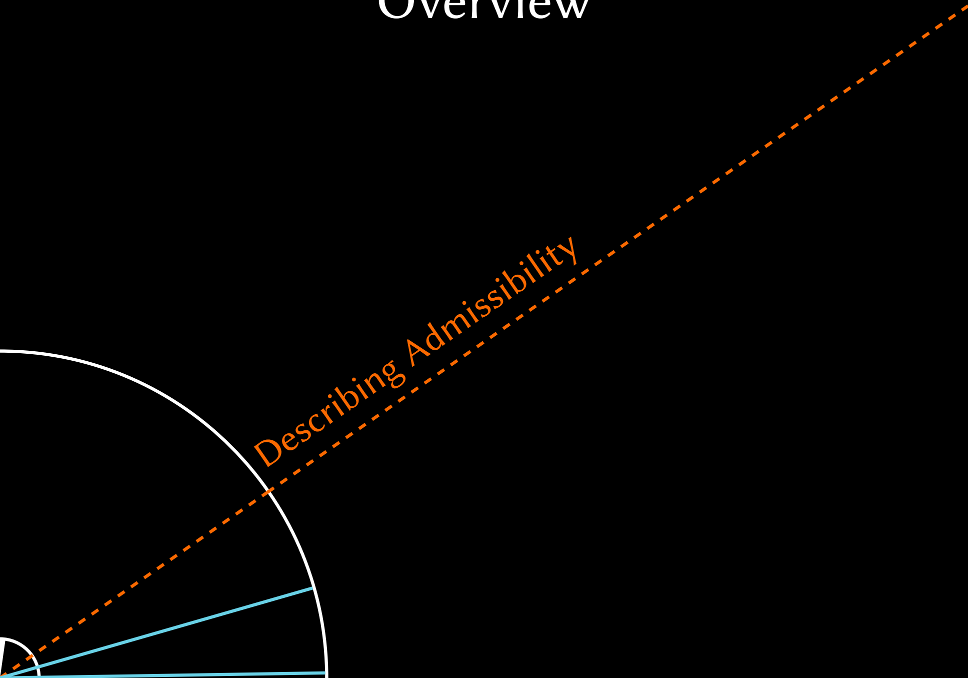

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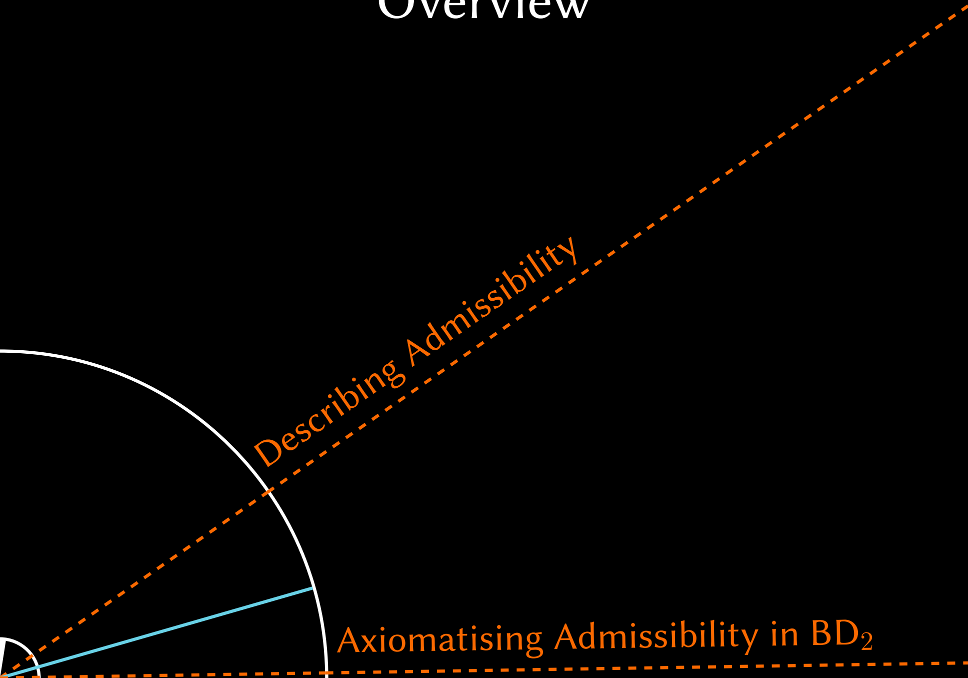
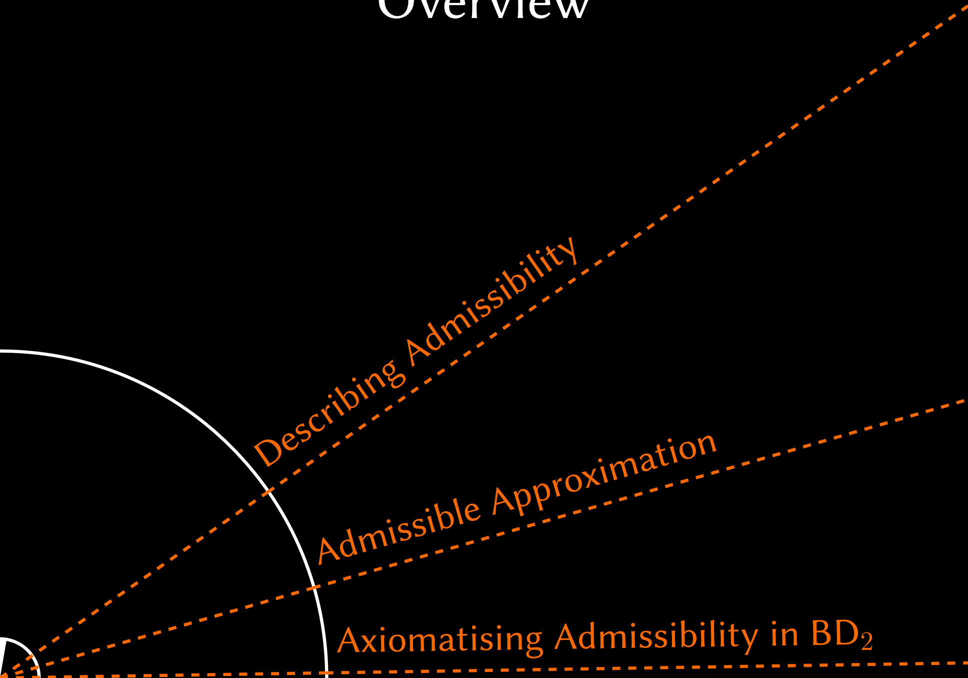


Describing Admissibility
Axiomatising Admissibility in BD_2

Overview


Describing Admissibility
Admissible Approximation
Axiomatising Admissibility in BD_2

A / $\triangle$ admissible

σA is derivable

$$A \;/\; \Delta \text{ admissible}$$

σC is derivable for some $C \in \Delta$

σA is derivable

$A \mathrel{|\!\sim} \Delta$ admissible

σC is derivable for some $C \in \Delta$

$$\frac{\neg C \to A \vee B}{(\neg C \to A) \vee (\neg C \to B)}$$

$$\frac{\neg C \to A \lor B}{\{\ \neg C \to A, \quad \neg C \to B\ \}}$$

1974 Gabbay and de Jongh

1974 Gabbay and de Jongh

1986 Maksimova

1952 Łukasiewicz

1974 Gabbay and de Jongh

1986 Maksimova

1952 Łukasiewicz

1957 Kreisel and Putnam

1974 Gabbay and de Jongh

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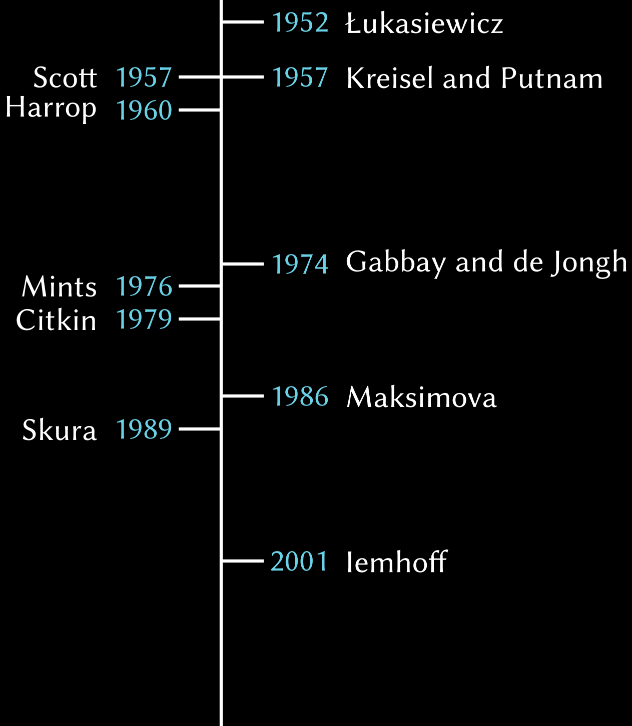
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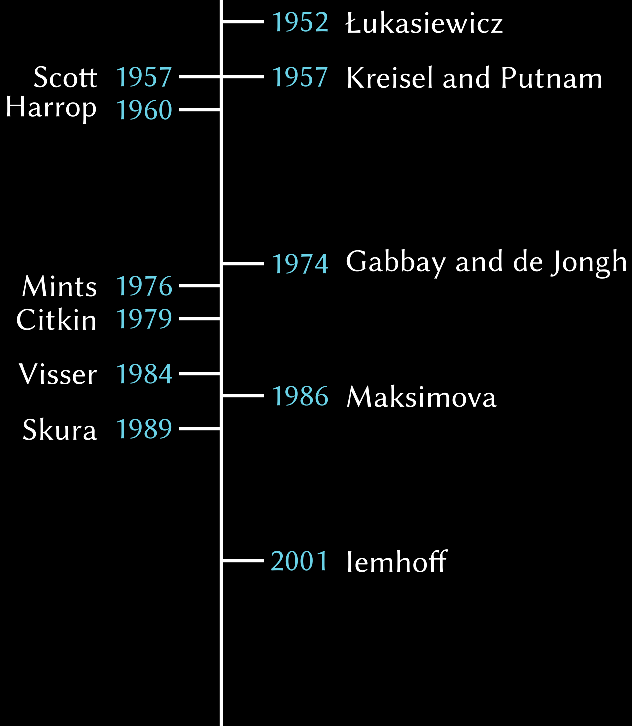
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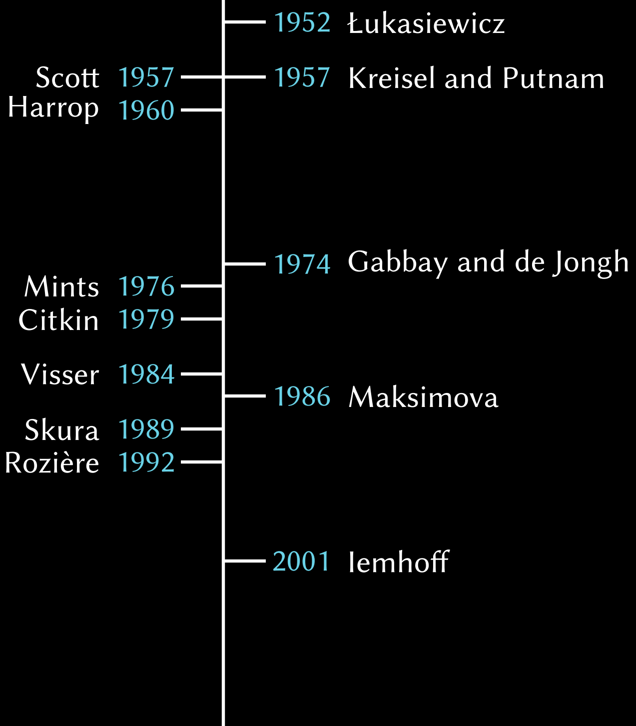
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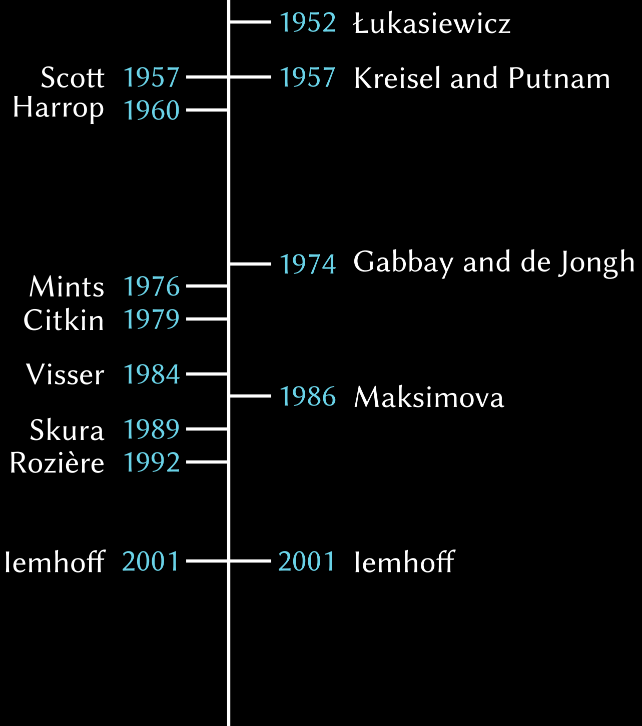
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G and Iemhoff 2014		

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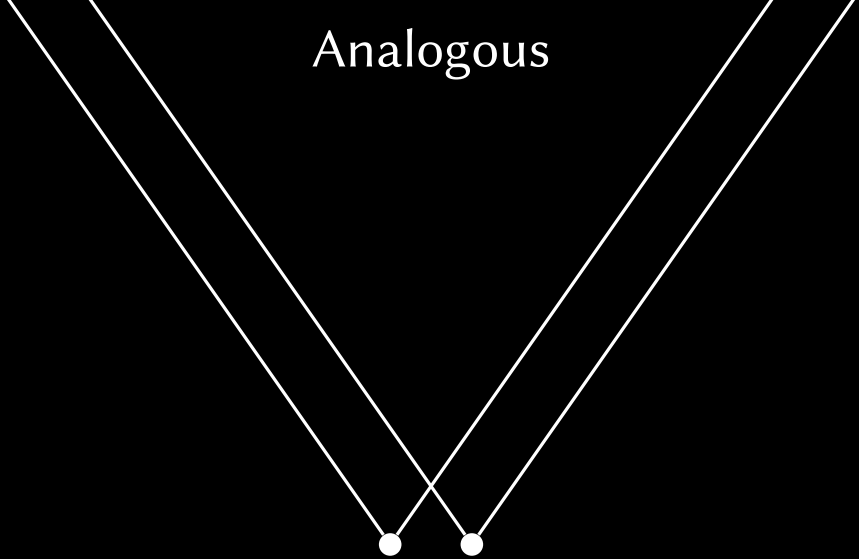
Mints	1976		
Citkin	1979		
Visser	1984		
		1986	Maksimova
Skura	1989		
Rozière	1992		
Iemhoff	2001	2001	Iemhoff
		2013	G
G and Iemhoff	2014		

Visser Rules

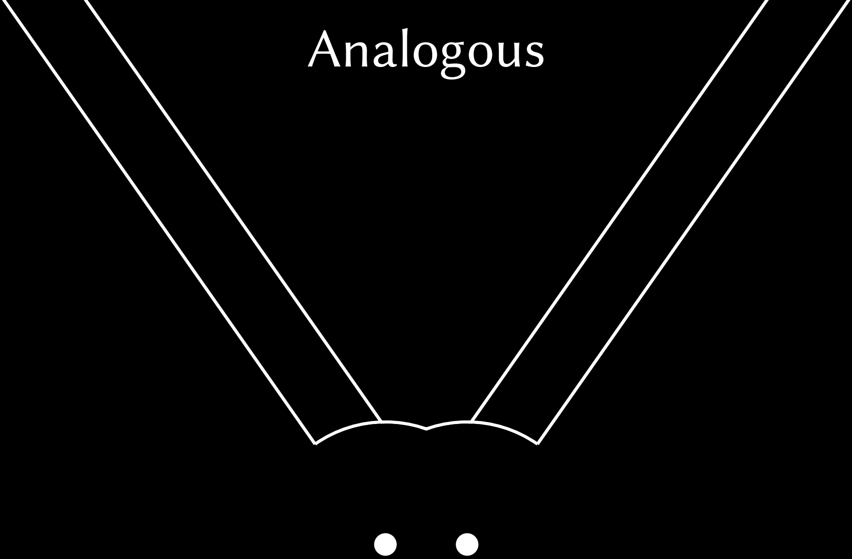
$$\frac{(\bigvee \Delta \to A) \to \bigvee \Delta}{\bigvee \left\{ (\bigvee \Delta \to A) \to C \;\middle|\; C \in \Delta \right\}}$$

Analogous

Analogous



Analogous



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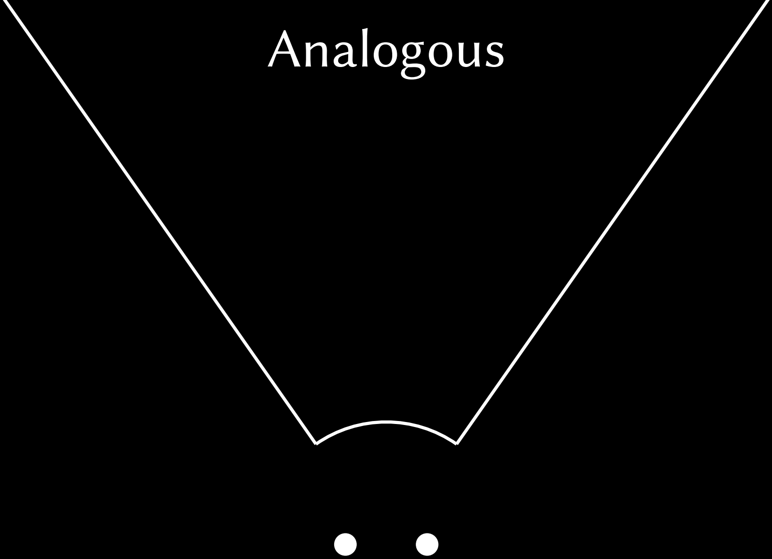

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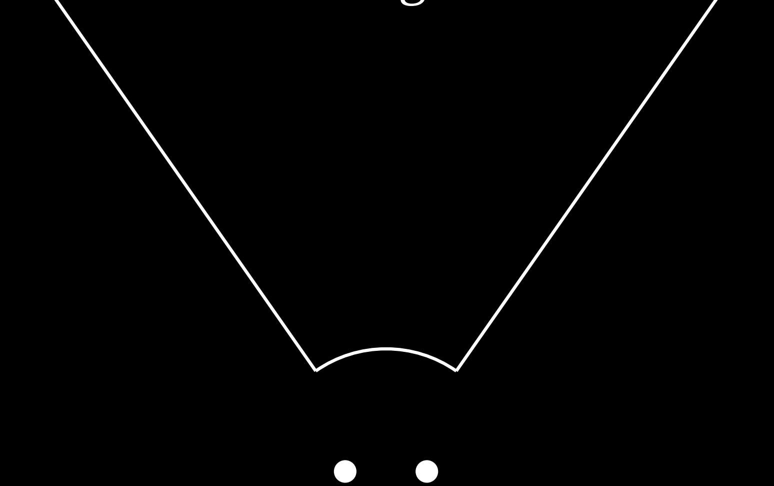
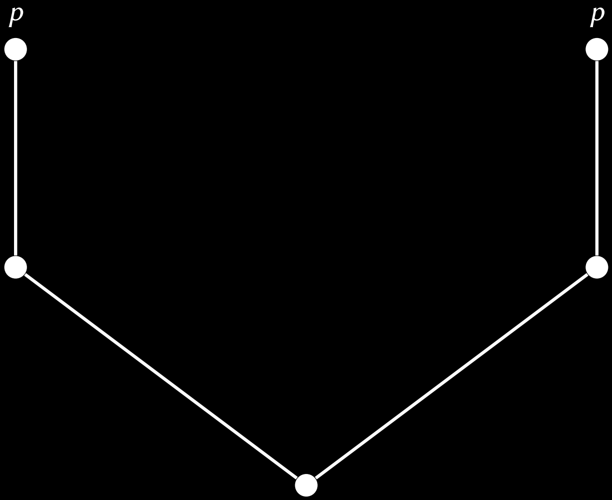


$k \equiv l$ when $v(k) = v(l)$ and $k \leq u$ iff $l \leq u$ for all $u \neq k, l$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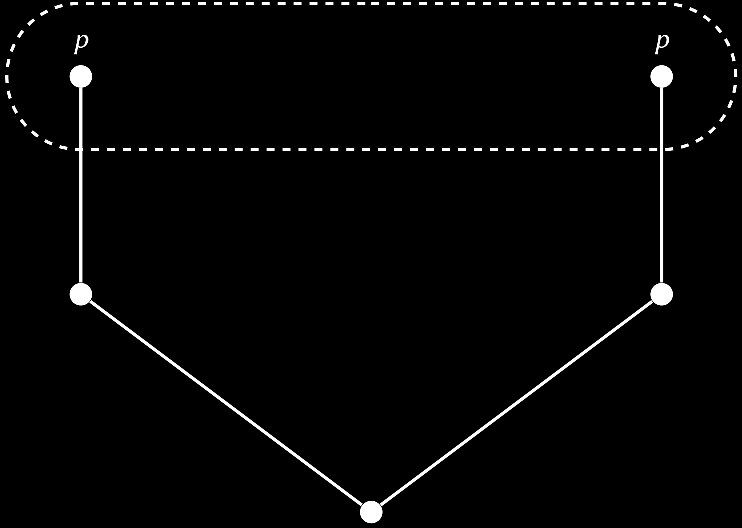
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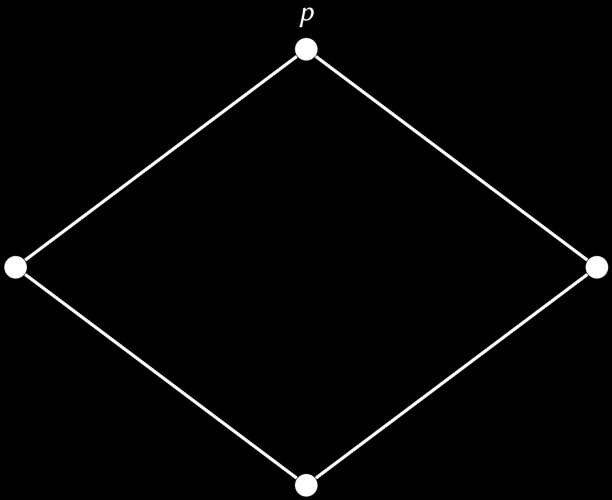
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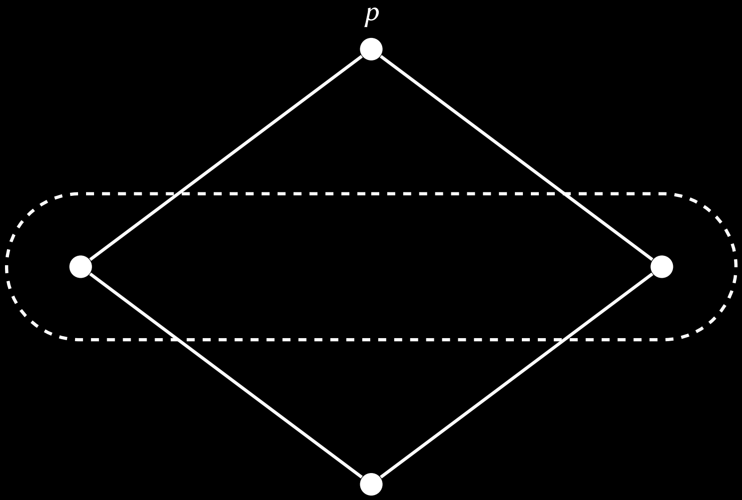
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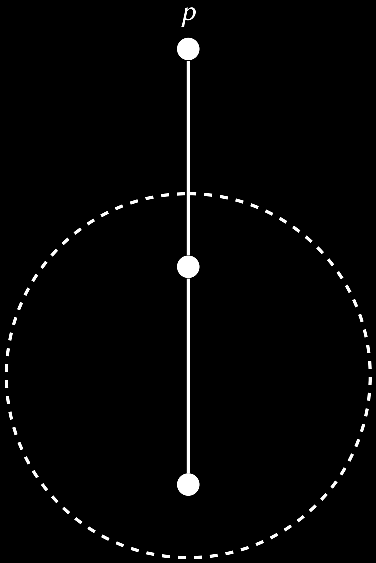
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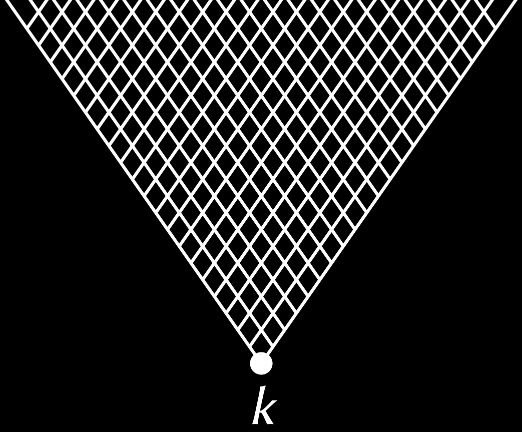
p

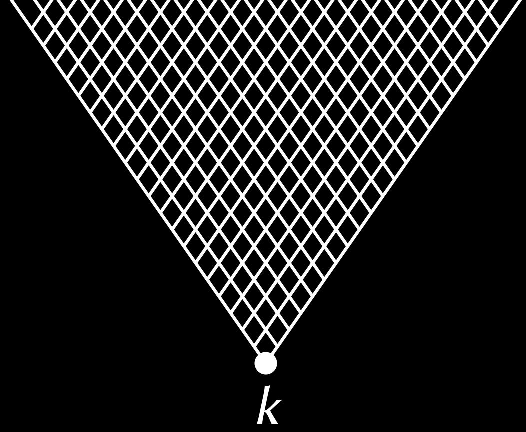
Jankov–de Jongh formulae

In suitable models have

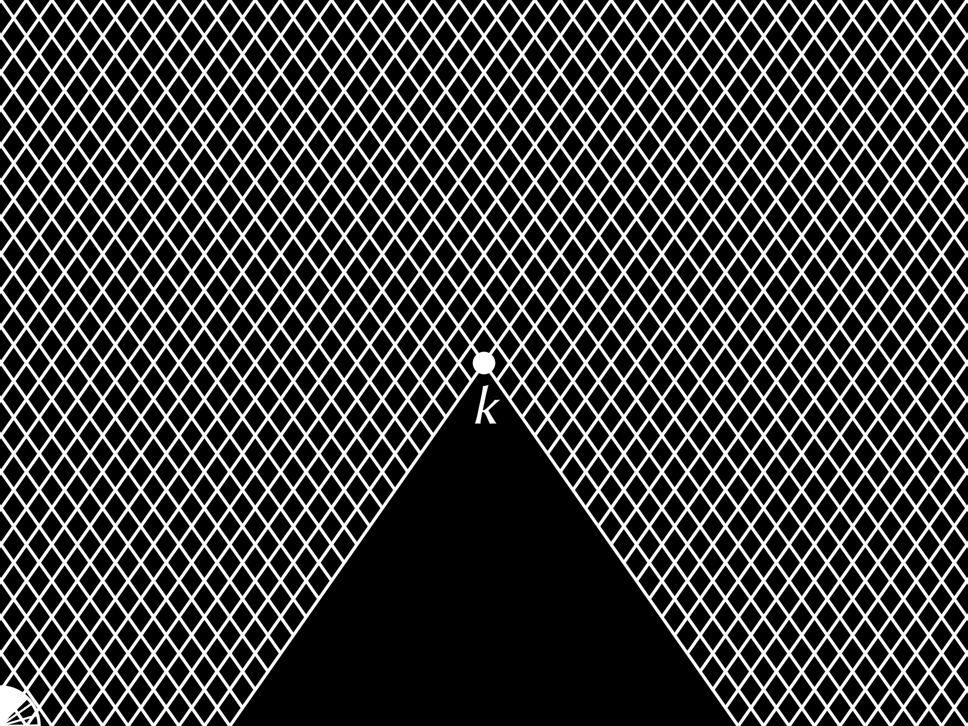
$$l \Vdash \mathrm{up}\, k \quad \text{iff} \quad k \leq l$$
$$l \Vdash \mathrm{nd}\, k \quad \text{iff} \quad l \not\leq k$$

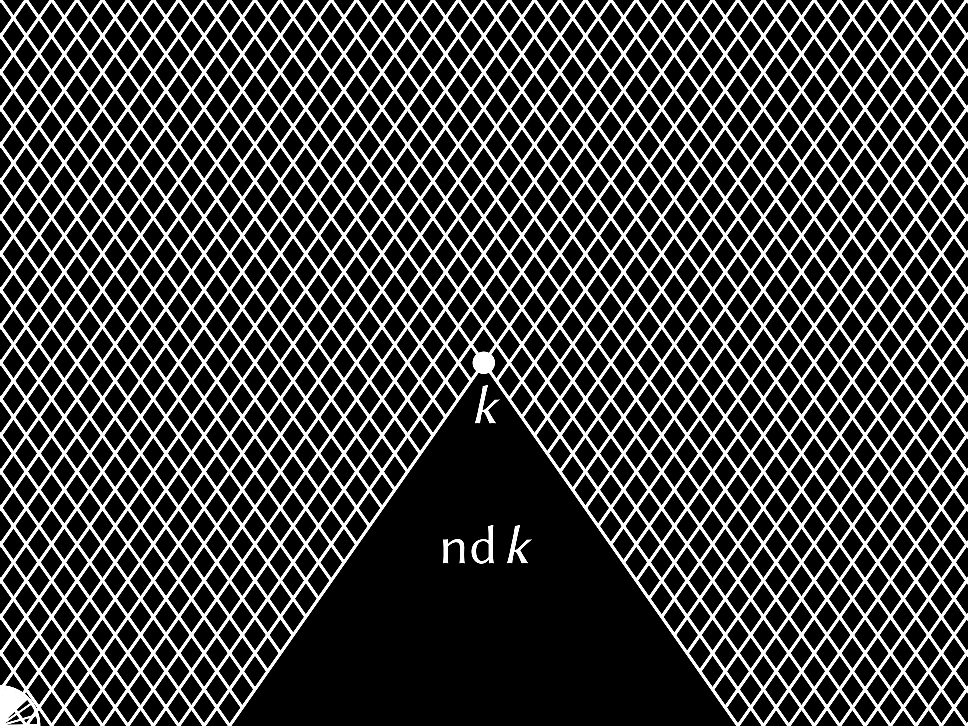

k

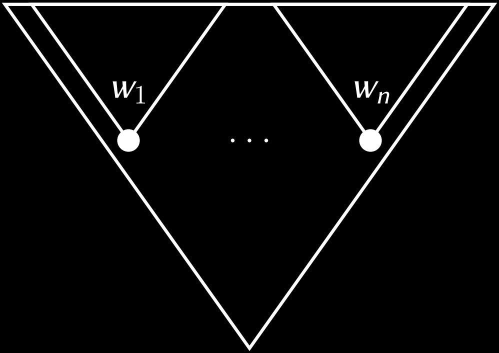

k



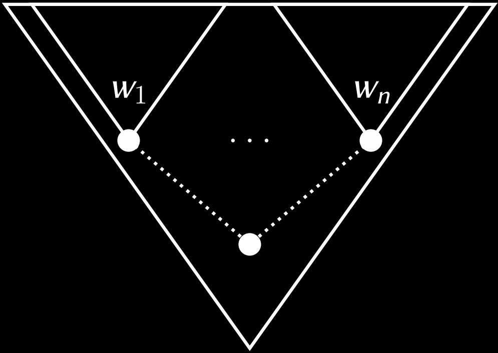
up k

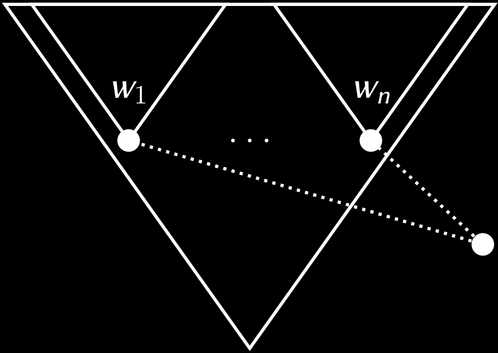

k


k
nd k

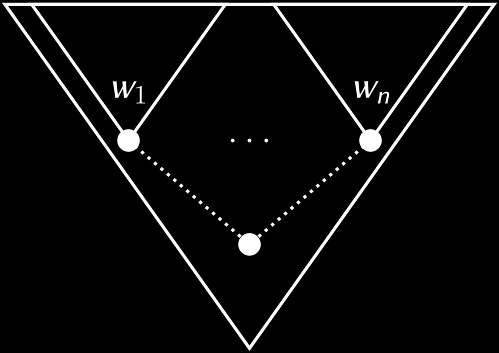


w_1
w_n
. . .

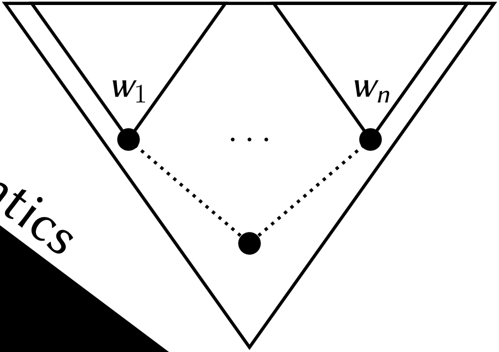

w_1
w_n
...

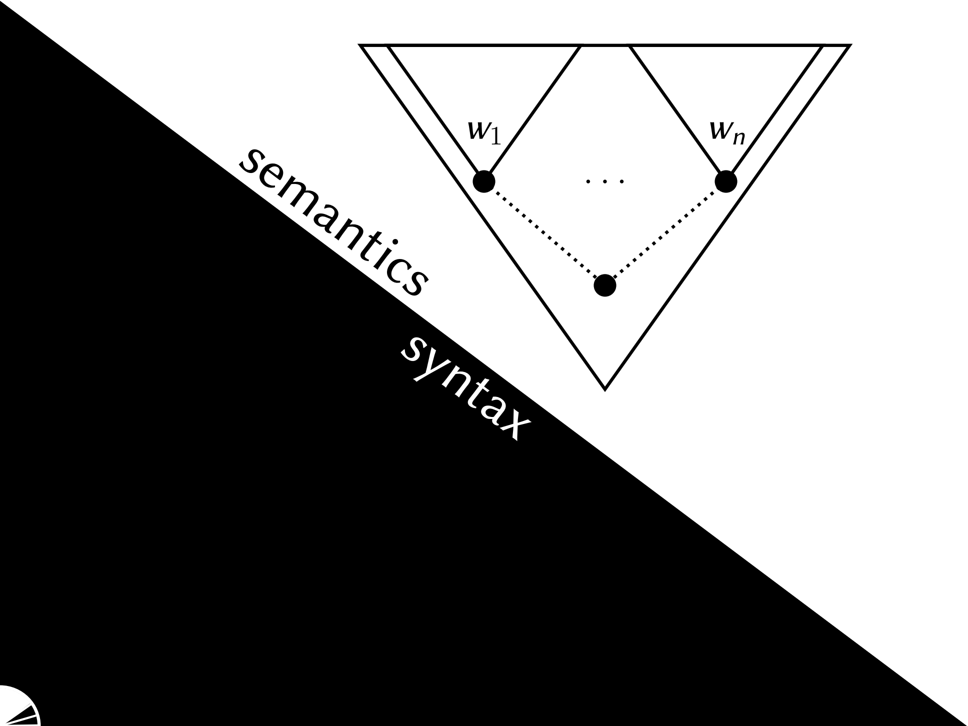


w_1
w_n


w_1
w_n
. . .

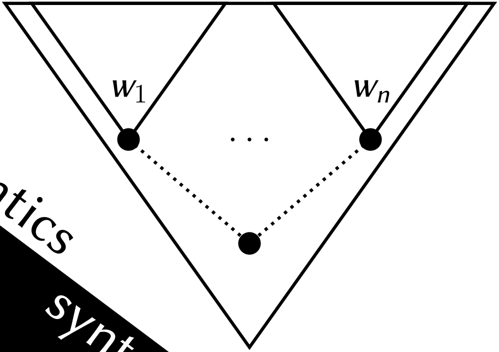
semantics



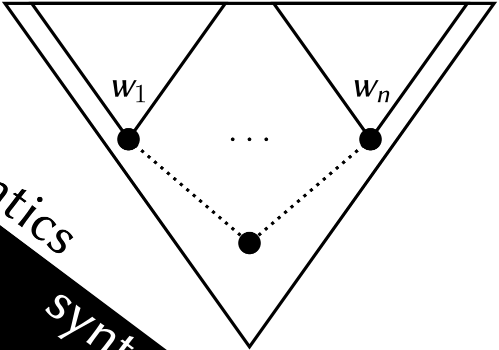

semantics
syntax
w_1
w_n
...

semantics

syntax



$$\frac{\left(\bigvee_{i=1}^{n} \mathsf{nd}\, w_i \rightarrow \bigvee_{i=1}^{n} \mathsf{up}\, w_i\right) \rightarrow \bigvee_{i=1}^{n} \mathsf{nd}\, w_i}{\bigvee_{j=1}^{n} \left(\bigvee_{i=1}^{n} \mathsf{nd}\, w_i \rightarrow \bigvee_{i=1}^{n} \mathsf{up}\, w_i\right) \rightarrow \mathsf{nd}\, w_j}$$



semantics

syntax

$$\frac{\left(\bigvee \Delta \quad \to A \qquad\right) \to \bigvee \Delta}{\bigvee_{C \in \Delta} \left(\bigvee \Delta \quad \to A \qquad\right) \to C}$$

An axiomatisation of admissibility is a set of rules R with

$$\vdash_R = \mathrel{|\!\!\sim}$$

Logic of Depth n

$$\begin{aligned} \mathrm{bd}_0 &= \bot \\ \mathrm{bd}_{n+1} &= p_{n+1} \vee (p_{n+1} \to \mathrm{bd}_n). \end{aligned}$$

Logic of Depth n

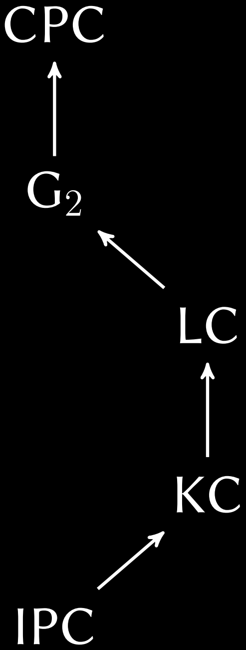
$$\begin{aligned} \mathrm{bd}_0 &= \bot \\ \mathrm{bd}_{n+1} &= p_{n+1} \vee (p_{n+1} \to \mathrm{bd}_n). \end{aligned}$$

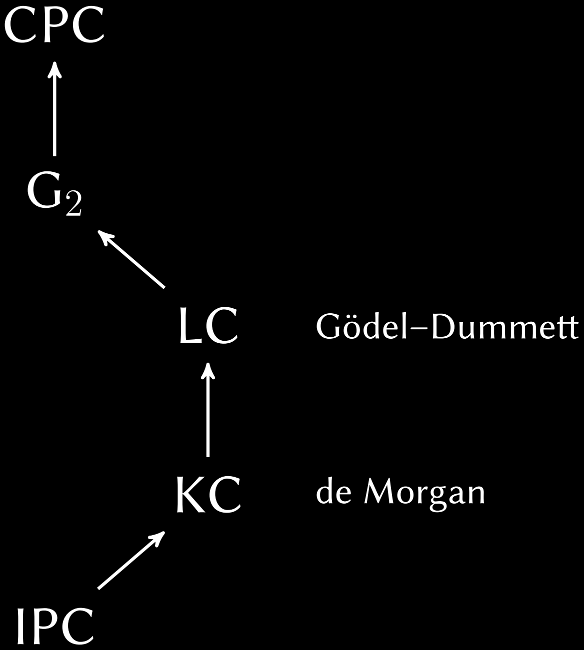
$$\mathrm{BD}_n = \mathrm{IPC} + \mathrm{bd}_n$$

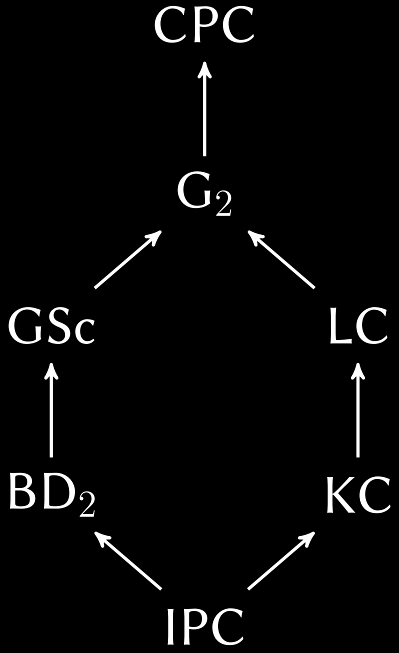
CPC

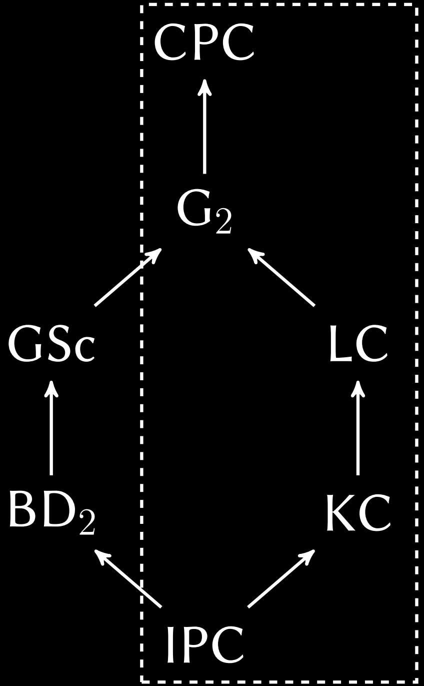
↑

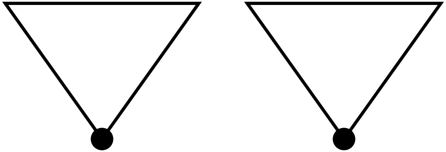
IPC

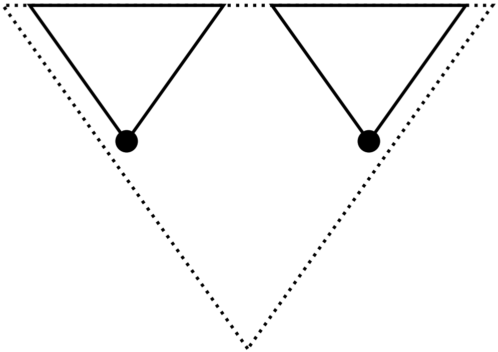

CPC
G2
LC
KC
IPC

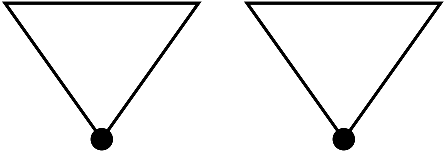

CPC
G_2
LC
Gödel–Dummett
KC
de Morgan
IPC

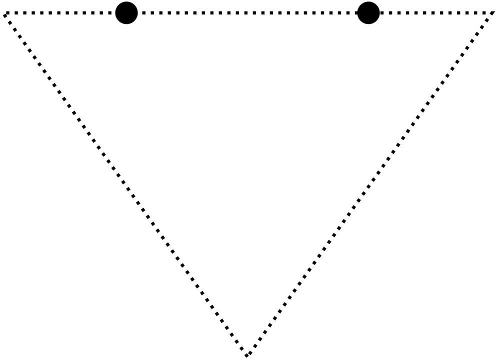

CPC
G_2
GSc
LC
BD_2
KC
IPC

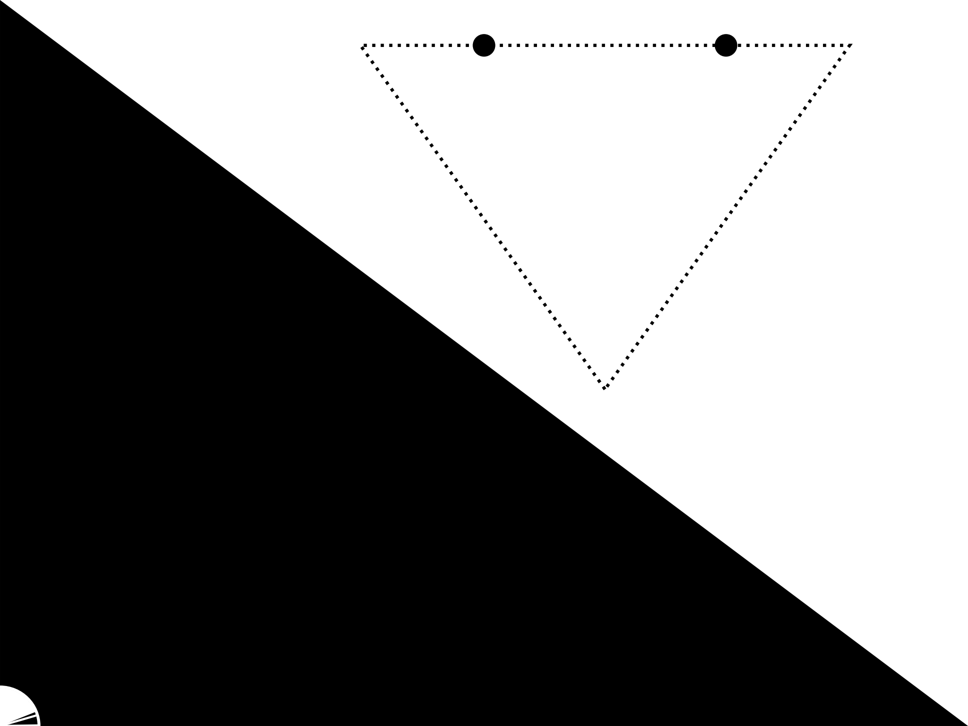

CPC
G_2
GSc
LC
BD_2
KC
IPC

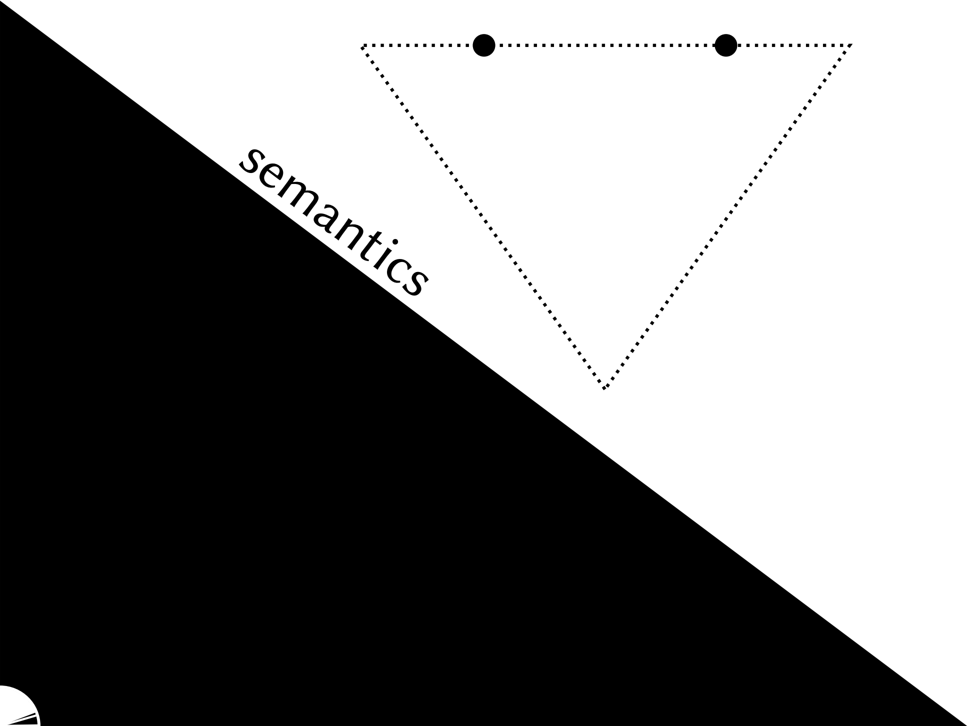








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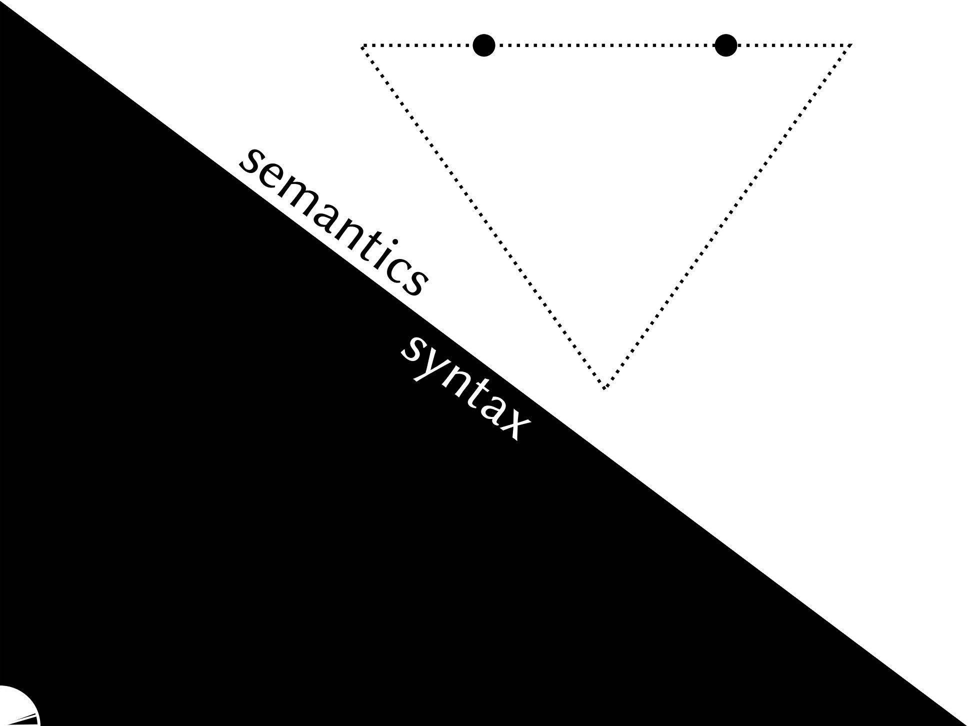
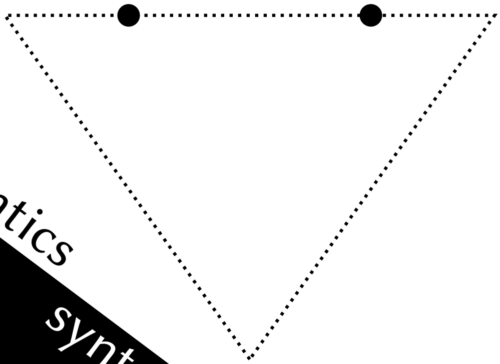

semantics
syntax



semantics

syntax

$$\frac{A \vee B}{\neg\neg A, \neg\neg B}$$

$$A \vdash B$$

$$\frac{A \vdash B}{A \mathrel{|\!\!\sim} B}$$

Admissible Approximation

$$\underline{A} \vdash B \text{ iff } A \mathrel{|}\!\!\sim B$$

If admissible approximations exists,
and if $A \vdash_{\mathrm{R}} \underline{A}$
then $\mid\!\sim\; \subseteq\; \vdash_{\mathrm{R}}$.

If admissible approximations exists,
and if $A \vdash_{\mathrm{R}} \underline{A}$ and $\mathrm{R} \subseteq \mathrel{|\!\sim}$
then $\mathrel{|\!\sim} = \vdash_{\mathrm{R}}$.

A is projective when

$\vdash \sigma A$ and $A \vdash \sigma B \equiv B$

for some σ.

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$\vdash \sigma A$ and $A \vdash \sigma B \equiv B$
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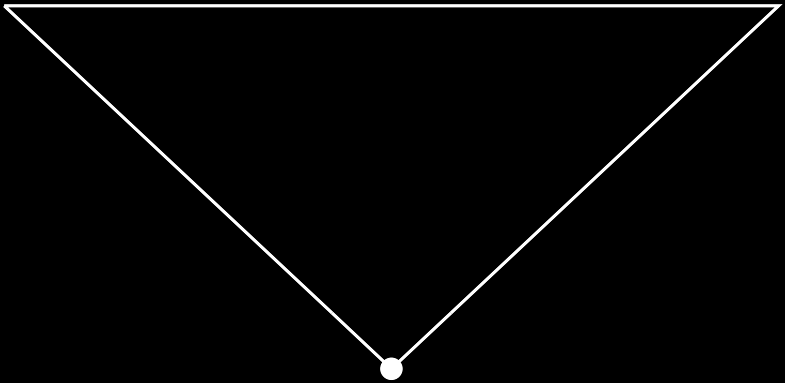
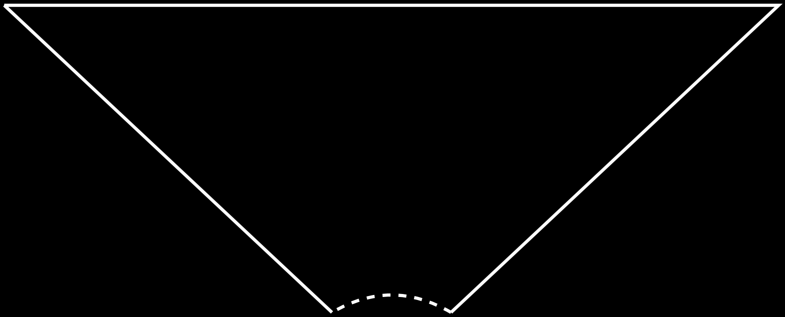
$$\underline{A} = A$$

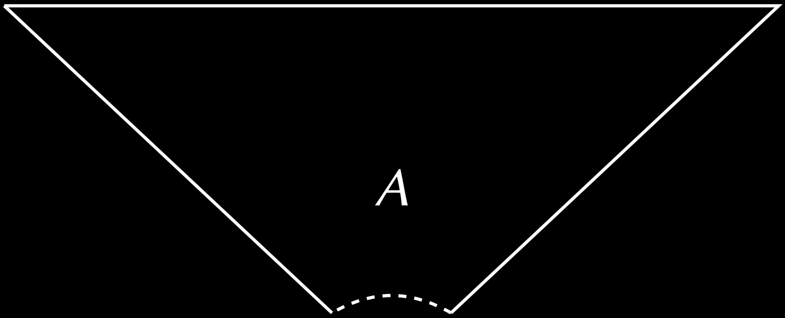
Ghilardi (1999)



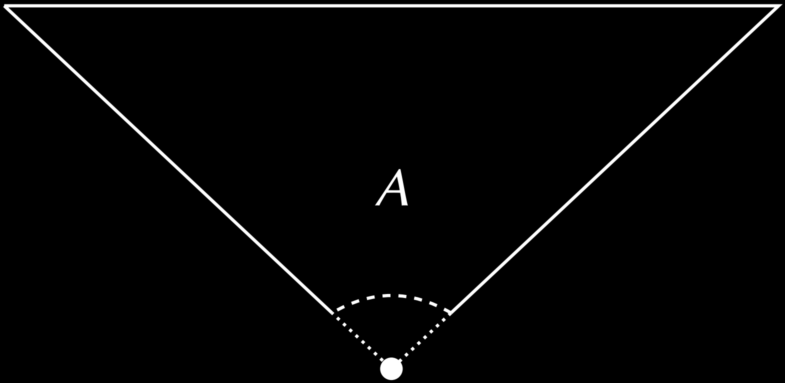
Ghilardi (1999)



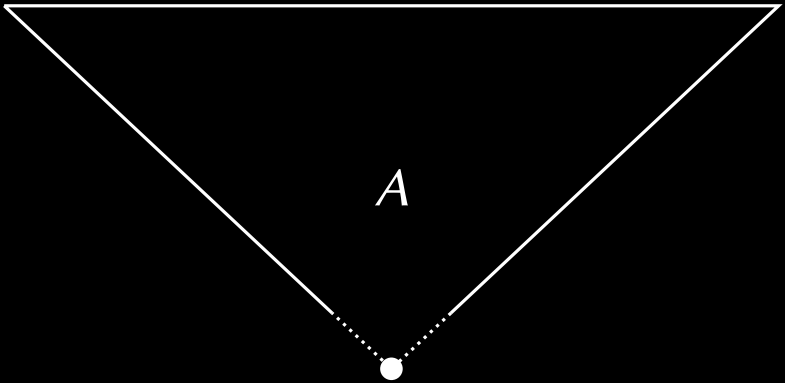
Ghilardi (1999)



Ghilardi (1999)



Ghilardi (1999)



Iemhoff (2001b)

A formula is IPC-projective iff it admits DP and V

Goudsmit and Iemhoff (2014)

A formula is T_n -projective iff
it admits DP and V_n
for $n \geq 2$

Visser Rules

$$\frac{(\bigvee \Delta \to A) \to \bigvee \Delta}{\bigvee \{ \quad (\bigvee \Delta \to A) \to C \mid C \in \Delta \}}$$

Skura (1992)

$$\frac{(\bigvee \Delta \to A) \to \bigvee \Delta}{\{\neg\neg((\bigvee \Delta \to A) \to C) \mid C \in \Delta\}}$$

A formula is BD_2-projective iff it admits S

To each A there is set Γ of BD_2-projectives with

$$A \vdash_S \bigvee \Gamma \text{ and } \bigvee \Gamma \vdash A$$

To each A there is set Γ of BD_2-projectives with

$$A \vdash_S \bigvee \Gamma \text{ and } \bigvee \Gamma \vdash A$$

which shows $\underline{A} = \bigvee \Gamma$.

Goudsmit (2013):

S axiomatises admissibility of BD_2

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