A Computational Approach to Finite MTL-chains

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Outline

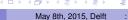
- Preliminaries
- Main Problem to Consider
- Reducing MTL-chains to involutive ones
- Representation Theorem for involutive MTL-chains
- 5 An Application: Exotic MTL-chains
- Final Remarks



A Really Quick Overview of the Framework

- language: \cdot , \vee , \wedge , 0, e, \rightarrow MTL: $\mathsf{FL}_{ew} + (x \rightarrow y) \lor (y \rightarrow x) = e$ (prelinearity)
 - ▶ BL: MTL + $x \land y = x \cdot (x \rightarrow y)$
 - ► IMTL: $MTL + (x \rightarrow 0) \rightarrow 0 = x$
 - ▶ IBL (MV): BL + $(x \rightarrow 0) \rightarrow 0 = x$





A Really Quick Overview of the Framework

language: ·, ∨, ∧, 0, e, →
 MTL: FL_{ew} + (x → y) ∨ (y → x) = e (prelinearity)
 BL: MTL + x ∧ y = x · (x → y)
 IMTL: MTL + (x → 0) → 0 = x

• ℓ -monoid (positive) language: $\cdot, \vee, \wedge, 0, e$

▶ IBL (MV): BL + $(x \rightarrow 0) \rightarrow 0 = x$



What are the ℓ-reducts of MTL-chains?

- \vee , \wedge is a linear partial order with bounds 0 and e,
- · is associative with neutral element *e* (i.e., monoid),
- · is commutative,
- · is monotone with respect to the partial order.



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When finite, w.l.o.g., it is enough to give the monoidal operation (under the assumption that the order is 0 < 1 < 2 ... < n).



Are these *ℓ*-reducts of MTL-chains?

	0	1	2	3	4
0	0	0	0	0	0
1	0	0	0	0	1
2	0	0	0	2	2
3	0	0	2	3	3
4	0	1	0 0 0 2 2	3	4

	0	1	2	3	4
0	0	0	0	0	0
1	0	0	0	0	1
2	0	0	1	2	2
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0	0	0	0	0	
0	0	0	0	1	
0	0	0	2	2	
0	0	2	3	3	
0	1	2	3	4	
	0 0 0 0 0	0 1 0 0 0 0 0 0 0 0 0 1	0 1 2 0 0 0 0 0 0 0 0 0 0 0 2 0 1 2	0 1 2 3 0 0 0 0 0 0 0 0 0 0 0 2 0 0 2 3 0 1 2 3	0 1 2 3 4 0 0 0 0 0 0 0 0 0 1 0 0 0 2 2 0 0 2 3 3 0 1 2 3 4

	0	1	2	3	4
0	0	0	0	0	0
1	0	0	0	0	1
2	0	0	1	2	2
3	0	0	2	3	3
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Associativity is the only non-trivial property to check.



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 (Nilpotent Minimum)



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An Embarrassing Question

Is the equation

$$x_1 x_4 x_7 \wedge x_2 x_5 x_8 \wedge x_3 x_6 x_9 \leq x_1 x_2 x_3 \vee x_4 x_5 x_6 \vee x_7 x_8 x_9$$



Is there some ℓ -monoid equation that distinguishes MTL from BL?

- No?
- Yes?

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Is the equation

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Is there some $\ell\text{-monoid}$ equation that distinguishes MTL from BL?

- No? (by HSP Theorem we would get a representation description for MTL-algebras)
- Yes? (requires a better understanding of (finite) MTL-chains)

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Is the equation

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Advertising slogan

Why considering the ℓ-monoid reduct?



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a better understanding of the ℓ-monoid fragment of MTL-algebras will enlighten us with a better understanding of the full language (including residuum).





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Why considering the ℓ-monoid reduct?

- a better understanding of the ℓ-monoid fragment of MTL-algebras will enlighten us with a better understanding of the full language (including residuum).
- ② in some contexts it is easier to deal with the ℓ -monoid fragment than with the full language.





"Balance" between Pros and Cons

- full language $(\cdot, \vee, \wedge, 0, e, \rightarrow)$
 - Pro:
 - Con:
- ℓ -monoid $(\cdot, \vee, \wedge, 0, e)$
 - Pro:
 - Con:





"Balance" between Pros and Cons

- full language $(\cdot, \vee, \wedge, 0, e, \rightarrow)$
 - Pro: Congruences are nicely characterized
 - Con:
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 - Con:





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 - Pro: Congruences are difficult
 - Con: Free abelian monoid is easy



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- full language $(\cdot, \vee, \wedge, 0, e, \rightarrow)$
 - Pro: Congruences are nicely characterized
 - Con: Free algebra is difficult
- ℓ-monoid (·, ∨, ∧, 0, e)
 - Pro: Congruences are difficult
 - ▶ Con: Free abelian monoid is easy, it is $\bigoplus_{i \in \kappa} (\mathbb{N}, +, 0)$







$$x \wedge y = y \wedge x$$
 $x \vee y = y \vee x$
 $x \wedge (y \wedge z) = (x \wedge y) \wedge z$ $x \vee (y \vee z) = (x \vee y) \vee z$
 $x \wedge (x \vee y) = x$ $x \vee (x \wedge y) = x$
 $x \wedge e = x$ $x \vee 0 = x$



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$$x \wedge (x \vee y) = x$$

$$x \wedge e = x$$

$$x \cdot (y \cdot z) = (x \cdot y) \cdot z$$

$$x \cdot y = y \cdot x$$

$$x \cdot e = x$$

$$x \cdot 0 = 0$$

$$x \lor y = y \lor x$$

$$x \lor (y \lor z) = (x \lor y) \lor z$$

$$x \lor (x \land y) = x$$

$$x \lor 0 = x$$



$$x \wedge y = y \wedge x \qquad x \vee y = y \vee x$$

$$x \wedge (y \wedge z) = (x \wedge y) \wedge z \qquad x \vee (y \vee z) = (x \vee y) \vee z$$

$$x \wedge (x \vee y) = x \qquad x \vee (x \wedge y) = x$$

$$x \wedge e = x \qquad x \vee 0 = x$$

$$x \cdot (y \cdot z) = (x \cdot y) \cdot z$$

$$x \cdot y = y \cdot x$$

$$x \cdot e = x$$

$$x \cdot 0 = 0$$

$$x \cdot (y \vee z) = (x \cdot y) \vee (x \cdot z) \qquad x \wedge (y \vee z) = (x \wedge y) \vee (x \wedge z)$$

$$x \cdot (y \wedge z) = (x \cdot y) \wedge (x \cdot z)$$

- A is a chain
- A has a coatom
- A is involutive



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$$\forall x(x \approx \neg \neg x)$$



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$$\forall xy(x < y \Rightarrow y \cdot \neg x \neq 0)$$

$$\forall xy(x < y \Rightarrow \exists z(x \cdot z = 0 \& y \cdot z \neq 0))$$



12



Involutive MTL-algebras

 The residuum operation can be replaced, in the signature, with negation because the equation

$$x \rightarrow y = \neg(x \cdot \neg y)$$

holds.





Involutive MTL-algebras

 The residuum operation can be replaced, in the signature, with negation because the equation

$$X \rightarrow Y = \neg (X \cdot \neg Y)$$

holds.

• Duality: if in an equation which is valid in an IMTL-chain \boldsymbol{A} and which only uses the symbols $\cdot, +, \neg, \wedge, \vee, 0, \boldsymbol{e}$ we simultaneously interchange

 \cdot and +, \wedge and \vee , \leq and \geq , 0 and e, then the resultant equation is also valid in the same A.





Remark (Duality)

For every IMTL-algebra **A**, the following two conditions are equivalent.

The equation

$$x_1 x_4 x_7 \wedge x_2 x_5 x_8 \wedge x_3 x_6 x_9 \leq x_1 x_2 x_3 \vee x_4 x_5 x_6 \vee x_7 x_8 x_9$$

is valid in A.

The equation

$$(x_1 + x_4 + x_7) \lor (x_2 + x_5 + x_8) \lor (x_3 + x_6 + x_9) \ge$$

 $\ge (x_1 + x_2 + x_3) \land (x_4 + x_5 + x_6) \land (x_7 + x_8 + x_9)$

is valid in A.



	1	2	3	4	5	6	7	8	9
BL	1	1	2	2 ²	2 ³	2^4	2 ⁵	2 ⁶	2 ⁷



	1	2	3	4	5	6	7	8	9
BL	1	1	2	2^2	2 ³	2^4	2 ⁵	2^{6}	2 ⁷
IBL (MV)	1	1	1	1	1	1	1	1	1



	1	2	3	4	5	6	7	8	9
BL	1	1	2	2 ²	2 ³	2^4	2 ⁵	2 ⁶	2 ⁷
IBL (MV)	1	1	1	1	1	1	1	1	1
MTL	1	1	2	6	22	94	451	2386	13775

A030453 Number of linearly ordered Abelian monoids of size n (semi-groups with 1 greatest element of the corresponding chain as neutral element): triangular norms on an n-chain. 1, 1, 2, 6, 22, 94, 451, 2386, 13775, 86417, 590489, 4446029, 37869449, 382549464 (list; graph; refs; listen; history; text; internal format) OFFSET 1.3 Table of n, a(n) for n=1..14. LINKS Bernard de Baets, Home page Index entries for sequences related to monoids CROSSREES Sequence in context: A150274 A109317 A109153 * A001861 A049526 A187251 Adjacent sequences: A030450 A030451 A030452 * A030454 A030455 A030456 KEYWORD nonn.hard.nice AUTHOR Bernard De Baets (Bernard.DeBaets(AT)rug.ac.be) STATUS approved





	1	2	3	4	5	6	7	8	9
BL	1	1	2	2 ²	2 ³	2^4	2 ⁵	2 ⁶	2 ⁷
IBL (MV)	1	1	1	1	1	1	1	1	1
MTL	1	1	2	6	22	94	451	2386	13775
IMTL	1	1	1	2	3	7	12	31	59

A034786 Number of linearly ordered Girard monoids of size n: number of t-norms of on an n-chain inducing an involutive residual negator. 1, 1, 1, 2, 3, 7, 12, 31, 59, 161, 329, 944, 2067, 6148, 14558, 44483, 116372 (list; graph; refs; listen: history: text; internal format) OFFSET M. Nachtegael. The Dizzy Number of Fuzzy Implication Operators on Finite Chains. REFERENCES in "Fuzzy Logic and Intelligent Technologies for Nuclear Science and Industry", ed. Ruan D., Abderrahim H., D'hondt P., Kerre E., 1998, pp. 29-35. Table of n, a(n) for n=1..17. LINKS Index entries for sequences related to monoids Cf. A030453. CROSSREFS Sequence in context: A134565 A100982 A186009 * A080107 A056156 A112837 Adjacent sequences: A034783 A034784 A034785 * A034787 A034788 A034789 KEYWORD Bernard De Baets (Bernard.DeBaets(AT)rug.ac.be), Mike Nachtegael AUTHOR (mike.nachtegael(AT)rug.ac.be) STATUS approved





*(5,0)	0	1	2	3	4	*(5, 1)	0	1	2	3	4	*(5,2)	0	1	2	3	4
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
1	0	0	0	0	1	1	0	0	0	0	1	1	0	0	0	0	1
2	0	0	0	2	2	2	0	0	0	1	2	2	0	0	0	1	2
3	0	0	2	3	3	3	0	0	1	2	3	3	0	0	1	1	3
4	0	1	2	3	4	4	0	1	2	3	4	4	0	1	2	3	4



*(5,0)	0	1	2	3	4	*(5, 1)	0	1	2	3	4	*(5, 2)	0	1	2	3	4
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
1	0	0	0	0	1	1						1					
2	0	0	0	2	2	2	0	0	0	1	2	2	0	0	0	1	2
3 4	0	0	2	3	3	$\frac{3}{4}$	0	0	1	2	3	3	0	0	1	1	3
4	0	1	2	3	4	4	0	1	2	3	4	4	0	1	2	3	4

+(5,0)	0	1	2	3	4	+(5,1)	0	1	2	3	4	+(5,2)	0	1	2	3	4
0	0	1	2	3	4	0	0	1	2	3	4	0	0	1	2	3	4
1	1	1	2	4	4	1	1	2	3	4	4	1	1	3	3	4	4
2	2	2	4	4	4	1 2 3 4	2	3	4	4	4	2	2	3	4	4	4
3	3	4	4	4	4	3	3	4	4	4	4	3	3	4	4	4	4
4	4	4	4	4	4	4	4	4	4	4	4	4	4	4	4	4	4



*(5,0	0)	0	1	2	3	4	*(5,1)	0	1	2	3	4		*(5, 2)	0	1	2	3	4
0		0	0	0	0	0	0	0	0	0	0	0	-	0	0	0	0	0	0
1		0	0	0	0	1	1							1	0	0	0	0	1
2	- 1	0	0	0	2	2	2	0	0	0	1	2		2	0	0	0	1	2
3		0	0	2	3	3	3	0	0	1	2	3		3	0	0	1	1	3
4		0	1	2	3	4	4	0	1	2	3	4		4	0	1	2	3	4

	+(5,0)	0	1	2	3	4	+(5,1)	0	1	2	3	4	+(5, 2)	0	1	2	3	4
Т	0	0	1	2	3	4	0	0	1	2	3	4	0	0	1	2	3	4
	1	1	1	2	4	4	1	1	2	3	4	4	1	1	3	3	4	4
	2	2	2	4	4	4	1 2 3 4	2	3	4	4	4	2	2	3	4	4	4
	3	3	4	4	4	4	3	3	4	4	4	4	3	3	4	4	4	4
	4	4	4	4	4	4	4	4	4	4	4	4	4	4	4	4	4	4

Strong Advice: Use the sage package created by Peter Jipsen



O	*(6,0)	0	1	2	3	4	5	*(6, 1)	0	1	2	3	4	5	*(6, 2)	0	1	2	3	4	5
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$		0	0	0	0	0	0		0	0	0	0	0	0	0	0	0	0	0	0	0
$\begin{array}{cccccccccccccccccccccccccccccccccccc$	1	0	0	0	0	0	1	1	0	0	0	0	0	1	1	0	0	0	0	0	1
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$\begin{array}{c ccccccccccccccccccccccccccccccccccc$	3	0	0	0	2	3	3	3	0	0	0	1	1	3	3	0	0	0	3	3	3
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$	4	0	0	2	3	4	4	4	0	0	1	1	1	4	4	0	0	2	3	4	4
0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0	5	0	1	2	3	4	5	5	0	1	2	3	4	5	5	0	1	2	3	4	5
0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0	(0.0)	La	_	_			_	(0.1)	La	_				_	/a =>	1	_	_			_
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$\begin{array}{c ccccccccccccccccccccccccccccccccccc$		0	0	0	0	2	2	2	0	0	0	0	1	2	2	0	0	0	0	1	2
5 0 1 2 3 4 5 5 0 1 2 3 4 5 5 0 1 2 3 4 5 5 0 1 2 3 4 5 6 6 6 6 6 6 6 6 6 6 6 6 6 6 6 6 6 6	3	0	0	0	1	2	3	3	0	0	0	3	3	3	3	0	0	0	1	1	3
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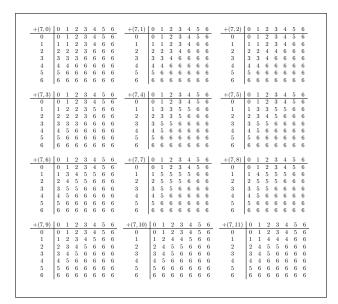
+(6,0)	0	1	2	3	4	5	+(6, 1)	0	1	2	3	4	5	+(6,2)	0	1	2	3	4	5
0	0	1	2	3	4	5	0	0	1	2	3	4	5	0	0	1	2	3	4	5
1	1	1	2	3	5	5	1	1	4	4	4	5	5	1	1	1	2	3	5	5
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3	3	3	5	5	5	5	3	3	4	5	5	5	5	3	3	3	5	5	5	5
4	4	5	5	5	5	5	4	4	5	5	5	5	5	4	4	5	5	5	5	5
5	5	5	5	5	5	5	5	5	5	5	5	5	5	5	5	5	5	5	5	5
+(6,3)	Lo	1	2	3	4	5	+(6,4)	10	1	2	3	4	5	+(6,5)	10	1	2	3	4	5
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1	1	1	3	3	5	5	1	1	2	2	4	5	5	1	1	3	4	4	5	5
2	2	3	4	5	5	5	2	2	2	2	5	5	5	2	2	4	4	5	5	5
3	3	3	5	5	5	5	3	3	4	5	5	5	5	3	3	4	5	5	5	5
4	4	5	5	5	5	5	4	4	5	5	5	5	5	4	4	5	5	5	5	5
5	5	5	5	5	5	5	5	5	5	5	5	5	5	5	5	5	5	5	5	5
. (0, 0)	Lo				,	-														
+(6,6)	-	1	2	3	4	5														
0	0	1	2	3	4	5														
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$ \begin{array}{cccccccccccccccccccccccccccccccccccc$	*(7, 3)	0	1	2	3	4	5	6	*(7, 4)	0	1	2	3	4	5	6	*(7,5)	0	1	2	3	4	5	6
$\begin{array}{cccccccccccccccccccccccccccccccccccc$	0	0	0	0	0	0	0	0	- 0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
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$\begin{array}{c ccccccccccccccccccccccccccccccccccc$	3	0	0	0	0	3	3	3	3	0	0	0	0	1	1	3	3	0	0	0	0	1	1	3
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$	4	0	0	0	3	4	4	4	4	0	0	0	1	3	3	4	4	0	0	0	1	2	3	4
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$	5	0	0	1	3	4	4	5	5	0	0	1	1	3	3	5	5	0	0	1	1	3	3	5
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$	6	0	1	2	3	4	5	6	6	0	1	2	3	4	5	6	6	0	1	2	3	4	5	6
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$ \begin{array}{c ccccccccccccccccccccccccccccccccccc$	6	0	1	2	3	4	5	6	6	0	1	2	3	4	5	6	6	0	1	2	3	4	5	6
$ \begin{array}{c ccccccccccccccccccccccccccccccccccc$	*(7.9)	Ιo	1	2	3	4	5	6	*(7.10)	Ιo	1	2	3	4	5	6	*(7.11)	Lο	1	2	3	4	5	6
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May 8th, 2015, Delft







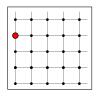
May 8th, 2015, Delft

Describing 2-generated IMTL-chains through Token Configurations in $(\mathbb{N}^2,+,0)$





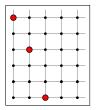




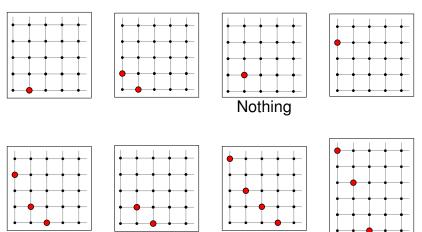






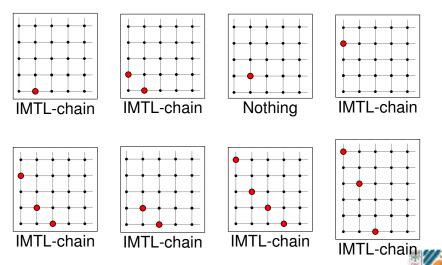


Describing 2-generated IMTL-chains through Token Configurations in $(\mathbb{N}^2, +, 0)$





Describing 2-generated IMTL-chains through Token Configurations in $(\mathbb{N}^2,+,0)$



antichain



• antichain (the "monomial ideal" generated is denoted I)



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- any of the following equivalent conditions hold:





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 - ▶ for all $a, b \in \mathbb{N}^2$, either $I \{a\} \subseteq I \{b\}$ or $I \{b\} \subseteq I \{a\}$





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 - ▶ for all $a, b, c, d \in \mathbb{N}^2$, if $a + b \in I$ and $c + d \in I$, then either $a + d \in I$ or $b + c \in I$
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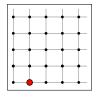


"Admissible" Configurations in $(\mathbb{N}^2, +, 0)$

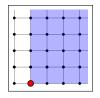
- antichain (the "monomial ideal" generated is denoted I)
- any of the following equivalent conditions hold:
 - $(\{I \{a\} : a \in \mathbb{N}^2\}, \subseteq)$ is a total order
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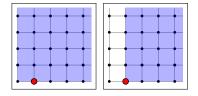
Communication Ideal: Any I with these conditions



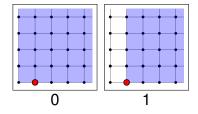




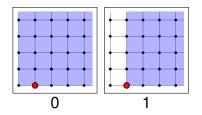






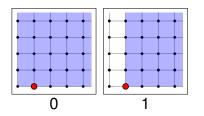






$$\begin{array}{c|cccc}
*(2,0) & 0 & 1 \\
\hline
0 & 0 & 0 \\
1 & 0 & 1
\end{array}$$

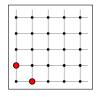




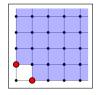
$$\begin{array}{c|cccc} *(2,0) & 0 & 1 \\ \hline 0 & 0 & 0 \\ 1 & 0 & 1 \\ \end{array}$$

2-element Boolean Algebra

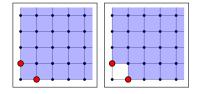




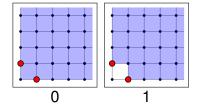




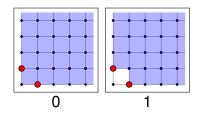






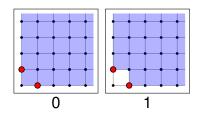






$$\begin{array}{c|cccc} *(2,0) & 0 & 1 \\ \hline 0 & 0 & 0 \\ 1 & 0 & 1 \\ \end{array}$$

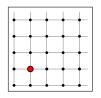




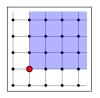
$$\begin{array}{c|cccc} *(2,0) & 0 & 1 \\ \hline 0 & 0 & 0 \\ 1 & 0 & 1 \\ \end{array}$$

2-element Boolean Algebra

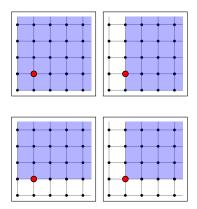




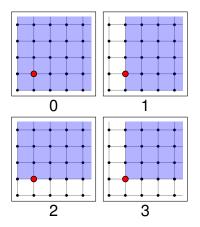




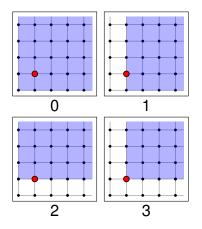












None (1 and 2 are not comparable)



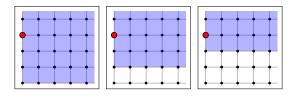
23





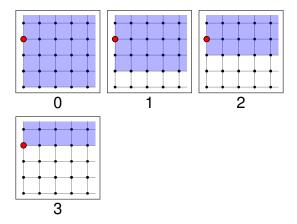




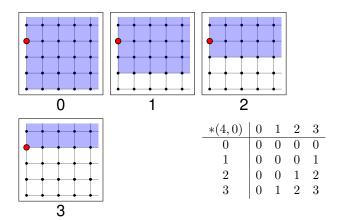




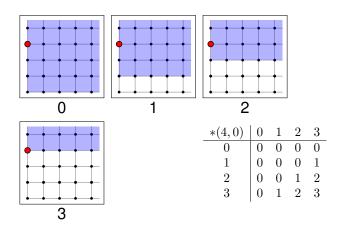








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4-element MV chain

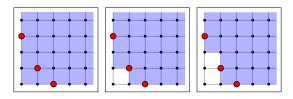






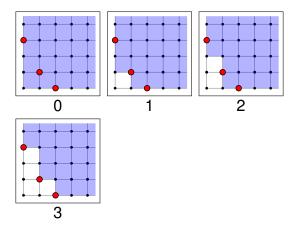




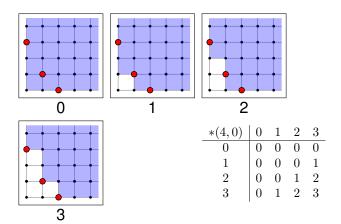




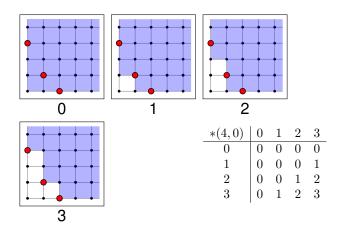






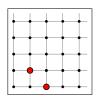




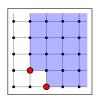


4-element MV chain

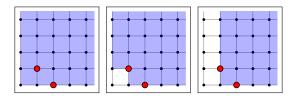






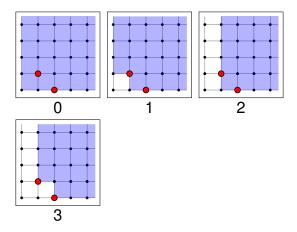




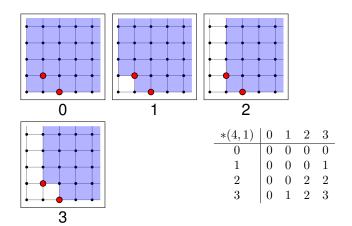




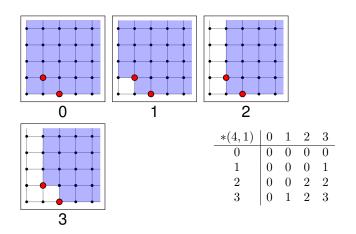








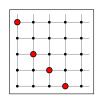




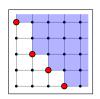
4-element Nilpotent Minimum chain



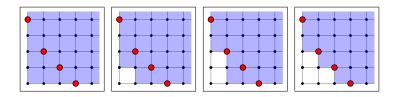
26

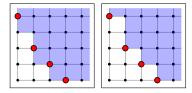




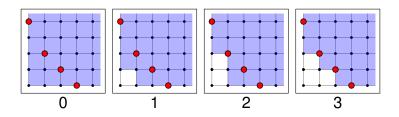


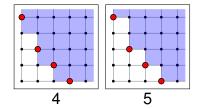




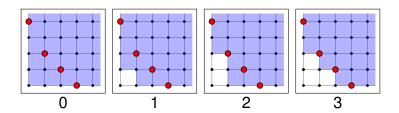


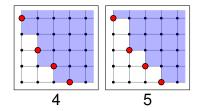






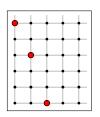




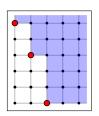


*(6,5)	0	1	2	3	4	5
0	0	0	0	0	0	0
1	0	0	0	0	0	1
2	0	0	0	0	1	2
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5	0	1	2	0 0 0 1 1 3	4	5

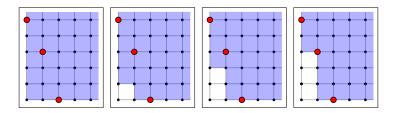


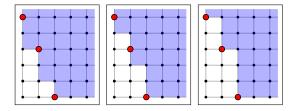




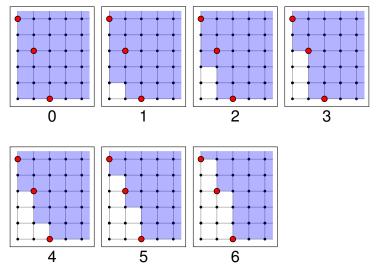




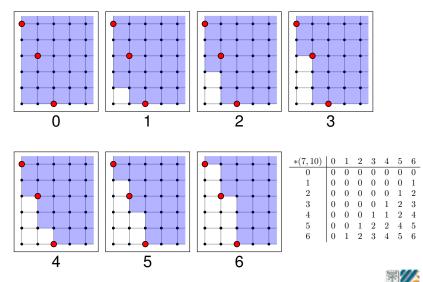












28

• Every algebra associated with a communication ideal of $(\mathbb{N},+,0)^2$ is a finite IMTL-chain.



- Every algebra associated with a communication ideal of $(\mathbb{N},+,0)^2$ is a finite IMTL-chain.
- All IMTL-chains that are 2-generated (as ℓ -monoid) come from a communication ideal of $(\mathbb{N},+,0)^2$





- Every algebra associated with a communication ideal of $(\mathbb{N},+,0)^2$ is a finite IMTL-chain.
- All IMTL-chains that are 2-generated (as ℓ -monoid) come from a communication ideal of $(\mathbb{N},+,0)^2$
- Every algebra associated with a communication ideal of $(\mathbb{N}, +, 0)^k$ is a finite IMTL-chain.



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- If a chain is *n*-potent, then we can replace the monoid $(\mathbb{N}, +, 0)$ with the "truncated" one over $\{0, 1, 2, \dots, n\}$.



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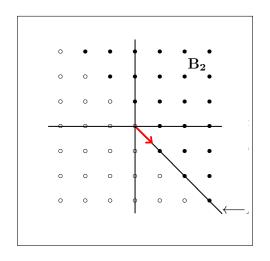
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- There are very nice geometrical interpretations of what are monomial orderings.



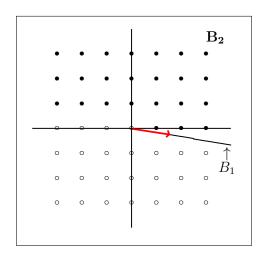
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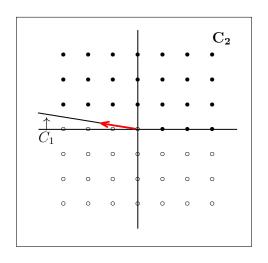


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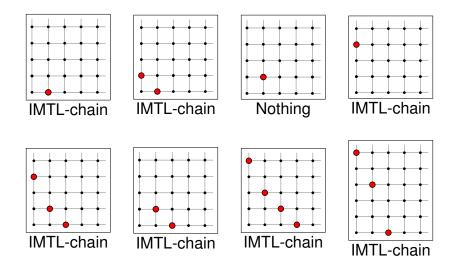


Two trivial ways to introduce communication ideals

- All upsets of admisible monomial orderings of $(\mathbb{N}^k, +, 0)$ are communication ideals.
 - For k = 2, communications ideals coincide exactly with (principal) upsets of admissible monomial orderings.
- The inverse image of a communication ideal under a monoid homomorphism is also a communication ideal.



Revisiting previous Token Configurations



The equation

 $X_1X_4X_7 \wedge X_2X_5X_8 \wedge X_3X_6X_9 \leq X_1X_2X_3 \vee X_4X_5X_6 \vee X_7X_8X_9$

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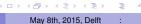
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- fails in MTL [Proof Sketch: explicit 36-element chain *E*]



```
#FUSION
fus table =
                                                        0.
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                                                           0.
                                                              0.
                                                              0,
                                 7,10,10,12,13,13,14,16,16,18,18,20,20,21,23,24,24,26,26,26,29,29,31,31,33,33,34],
                                 9,10,11,12,13,14,15,16,17,18,19,20,21,22,23,24,25,26,27,28,29,30,31,32,33,34,35],
```



Exotic MTL-chain:



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The algebra *E* is an exotic MTL-chain of dimension 9; the set of irreducible elements is {9, 15, 17, 22, 25, 28, 30, 32, 34}.





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This is the unique counterexample up to symmetry (and so there are 36 counterexamples)



Félix Bou (IIIA - CSIC) Finite MTL-chains May 8th, 2015, Delft :

```
#RESTDUUM
res table
              8. 9.10.11.12.13.14.15.16.17.18.19.20.21.22.23.24.25.26.27.28.29.30.31.32.33.34.35
 5
                                                          8
 / 31
                                                         1 32
               9.11.11.12.15.15.15.15.17.17.19.19.22.22.22.23.25.25.28.28.28.28.30.30.32.32.35.35.351.
             8, 9,10,11,12,13,14,15,16,17,18,19,20,21,22,23,24,25,26,27,28,29,30,31,32,33,34,35],
```



```
#ADDITION
add table
     8. 9.10.11.12.13.14.15.16.17.18.19.20.21.22.23.24.25.26.27.28.29.30.31.32.33.34.35
                     5
                     8
      / 31
                     1 32
```



How to obtain the previous exotic chain *E*?

Remember we want to falsify the equation

$$(x_1 + x_4 + x_7) \lor (x_2 + x_5 + x_8) \lor (x_3 + x_6 + x_9) \ge$$

 $\ge (x_1 + x_2 + x_3) \land (x_4 + x_5 + x_6) \land (x_7 + x_8 + x_9)$

 \bullet There is a communication ideal \emph{I} of $(\mathbb{N}^9,+,0)$ such that

$$e_1 + e_4 + e_7 \not\in I$$
, $e_2 + e_5 + e_8 \not\in I$, $e_3 + e_6 + e_9 \not\in I$
 $e_1 + e_2 + e_3 \in I$, $e_4 + e_5 + e_6 \in I$, $e_7 + e_8 + e_9 \in I$



$$h(a_1,a_2,\ldots,a_9) \coloneqq egin{pmatrix} 1 & 0 & 0 & 0 & 0 & 1 & 0 & 1 & 0 \ 0 & 1 & 0 & 0 & 0 & 1 & 1 & 0 & 0 \ 0 & 0 & 0 & 1 & 0 & -1 & -1 & 0 & 1 \ 0 & 0 & 1 & 0 & 0 & 1 & 1 & 1 & -1 \ 0 & 0 & 0 & 1 & -1 & 0 & -1 & 1 \end{pmatrix} egin{pmatrix} a_1 \ a_2 \ a_3 \ a_4 \ a_5 \ a_6 \ a_7 \ a_8 \ a_9 \end{pmatrix}$$

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• h is length-preserving and

$$h(e_1 + e_2 + e_3) = h(e_4 + e_5 + e_6) = h(e_7 + e_8 + e_9) = h(e_1 + e_4 + e_7) = h(e_2 + e_5 + e_8) = h(e_3 + e_6 + e_9) = h(e_1 + e_4 + e_7)$$



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- Claim: I is a communication ideal of $(\mathbb{N}, +, 0)^9$ satisfying our requirements.

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- Open: Is there some "very expressive" language that cannot distinguish MTL from BL? What about ·, ∨, 0, 1?



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involutive MTL-chains are "locally finite".



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- upsets of admissible monomial orderings provide an easy method to obtain communication ideals.
- an small perturbation method has been used to obtain a quite pathological example of communication ideal (its associated ITML-chain is exotic).

